

Metrics on diagrams and persistent homology

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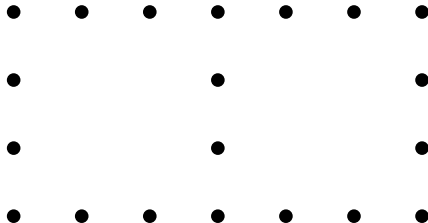
joint work with Vin de Silva and Jonathan A. Scott

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Topological data analysis

From data to topology:

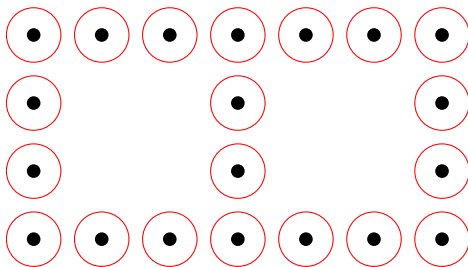
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- 2 Apply a geometric construction (e.g. Čech, Rips) to obtain a nested sequence of simplicial complexes.



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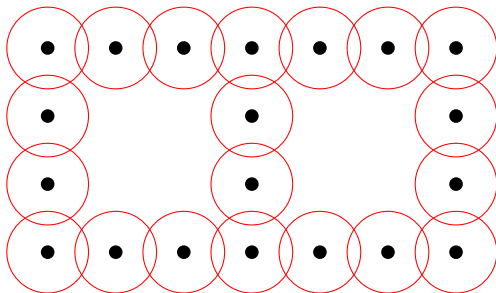
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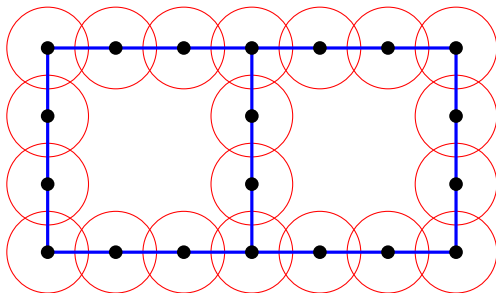
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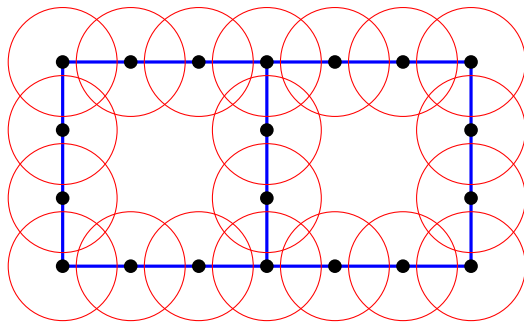
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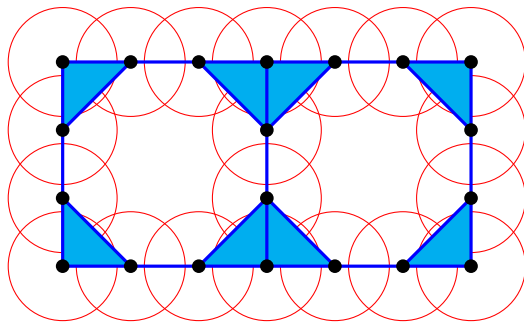
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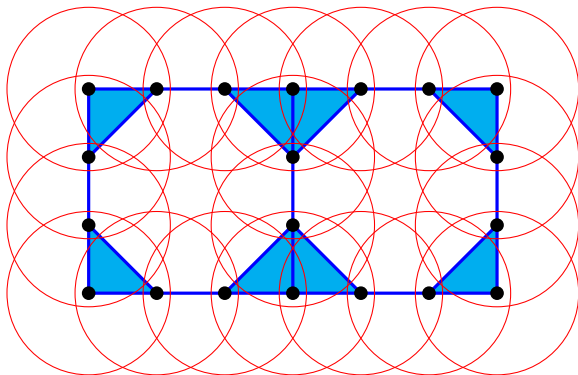
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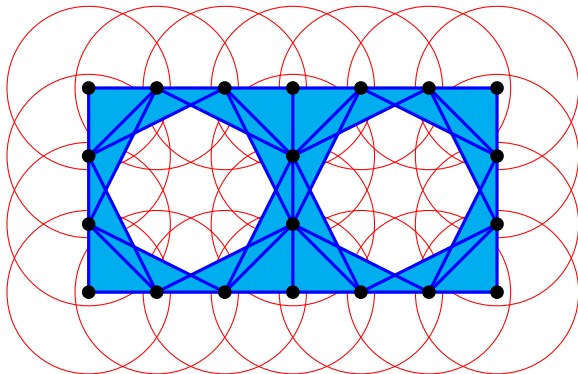
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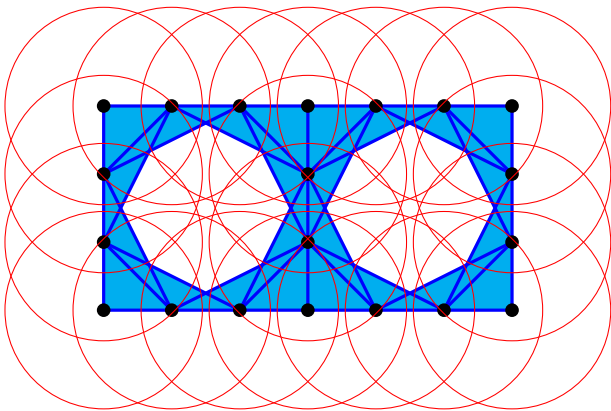
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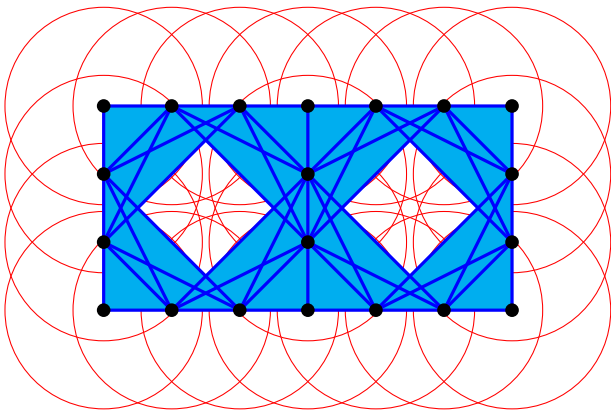
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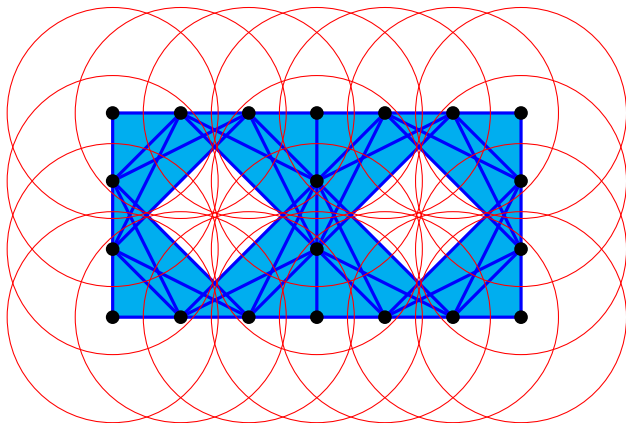
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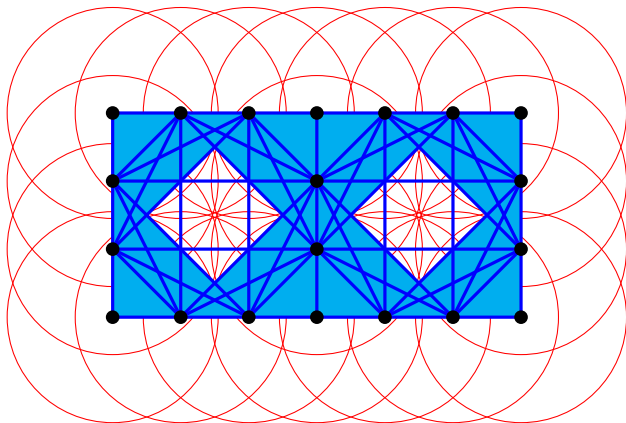
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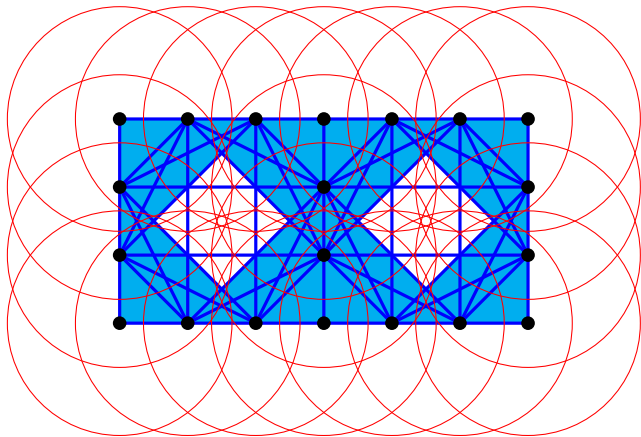
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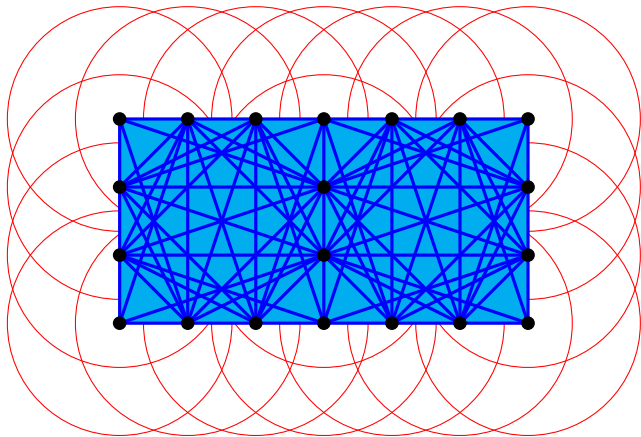
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Persistent homology

We have a nested sequence of simplicial complexes,

$$K_0 \longrightarrow K_1 \longrightarrow \cdots \longrightarrow K_n. \quad (*)$$

Apply simplicial homology,

$$H(K_0) \longrightarrow H(K_1) \longrightarrow \cdots \longrightarrow H(K_n). \quad (H^*)$$

Persistent homology

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Apply simplicial homology,

$$H(K_0) \longrightarrow H(K_1) \longrightarrow \cdots \longrightarrow H(K_n). \quad (H*)$$

The shape of these diagrams is given by the category \mathbf{n} ,

$$0 \longrightarrow 1 \longrightarrow \cdots \longrightarrow n.$$

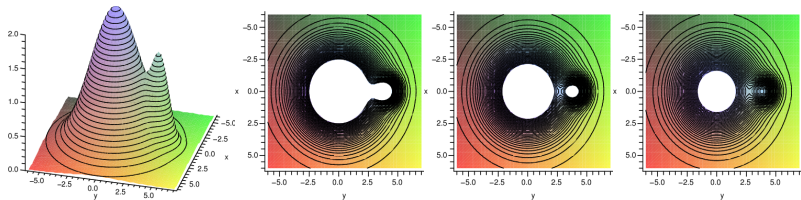
Then $(*)$ is equivalent to $\mathbf{n} \xrightarrow{K} \mathbf{Simp}$,

and $(H*)$ is equivalent to $\mathbf{n} \xrightarrow{K} \mathbf{Simp} \xrightarrow{H} \mathbf{Vect}_{\mathbb{F}}$.

Persistent homology

Another paradigm:

- 1 Start with a function $f : X \rightarrow \mathbb{R}$.
- 2
 - For each $a \in \mathbb{R}$, consider $f^{-1}((-\infty, a])$.
 - This gives us a diagram $F : (\mathbb{R}, \leq) \rightarrow \mathbf{Top}$.



- 3 Composing with singular homology we have,

$$(\mathbb{R}, \leq) \xrightarrow{F} \mathbf{Top} \xrightarrow{H} \mathbf{Vect}_{\mathbb{F}}$$

Multidimensional persistent homology

- 1 Start with a function $f : X \rightarrow \mathbb{R}^n$.
- 2
 - For each $a \in \mathbb{R}^n$, consider $f^{-1}(\mathbb{R}_{\leq a}^n)$.
 - This gives us a diagram $F : (\mathbb{R}^n, \leq) \rightarrow \mathbf{Top}$.
- 3 Composing with singular homology we have,

$$(\mathbb{R}^n, \leq) \xrightarrow{F} \mathbf{Top} \xrightarrow{H} \mathbf{Vect}_{\mathbb{F}}.$$

Levelset persistent homology

- 1 Start with a function $f : X \rightarrow \mathbb{R}$.
- 2
 - For each interval $I \subseteq \mathbb{R}$, consider $f^{-1}(I)$.
 - This gives us a diagram $F : \mathbf{Intervals} \rightarrow \mathbf{Top}$.
- 3 Composing with singular homology we have,

$$\mathbf{Intervals} \xrightarrow{F} \mathbf{Top} \xrightarrow{H} \mathbf{Vect}_{\mathbb{F}}$$

Angle-valued persistent homology

- 1 Start with a function $f : X \rightarrow S^1$.
- 2
 - For each arc $A \subseteq S^1$, consider $f^{-1}(A)$.
 - This gives us a diagram $F : \mathbf{Arcs} \rightarrow \mathbf{Top}$.
- 3 Composing with singular homology we have,

$$\mathbf{Arcs} \xrightarrow{F} \mathbf{Top} \xrightarrow{H} \mathbf{Vect}_{\mathbb{F}}.$$

Goals

We will use category theory to give a unified treatment of each of the above flavors of persistent homology.

Why?

- Give simpler, common proofs to some basic persistence results.
- Remove assumptions.
- Apply persistence to functions, $f : X \rightarrow (M, d)$.
- Allow homology to be replaced with other functors.
- Provide a framework for new applications.

Specific goal:

- Interpret and prove stability in this setting.

Terminology

In this talk a **metric** will be allowed to

- have $d(x, y) = \infty$ for $x \neq y$, and
- have $d(x, y) = 0$ for $x \neq y$.

That is, it is an extended pseudometric.

Example: The Hausdorff distance on the set of all subspaces of \mathbb{R} .

Unified framework

Generalized persistence module,

$$\mathbf{P} \xrightarrow{F} \mathbf{C} \xrightarrow{H} \mathbf{A}.$$

Here,

- The indexing category \mathbf{P} is a poset together with some notion of distance;
- \mathbf{C} is some category;
- \mathbf{A} is some abelian category (e.g. $\mathbf{Vect}_{\mathbb{F}}$, $\mathbf{R-mod}$);
- F and H are arbitrary functors.

Main results

Theorem (Interleaving distance)

There is a distance function $d(F, G)$ between diagrams $F, G : \mathbf{P} \rightarrow \mathbf{C}$.

This '*interleaving distance*' is a metric.

Theorem (Stability of interleaving distance)

Let $F, G : \mathbf{P} \rightarrow \mathbf{C}$ and $H : \mathbf{C} \rightarrow \mathbf{A}$. Then,

$$d(H \circ F, H \circ G) \leq d(F, G).$$

Inverse images of metric space valued functions

Start with $f : X \rightarrow (M, d_M)$.

Let \mathbf{P} be a poset of subsets of (M, d_M) .

Define

$$F : \mathbf{P} \rightarrow \mathbf{Top}$$

$$U \mapsto f^{-1}(U)$$

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Define

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$$U \mapsto f^{-1}(U)$$

Theorem (Inverse-image stability)

Let $F, G : \mathbf{P} \rightarrow \mathbf{Top}$ be given by inverse images of $f, g : X \rightarrow (M, d_M)$.

$$d(F, G) \leq d_\infty(f, g) := \sup_{x \in X} d_M(f(x), g(x)).$$

Corollary (Stability of generalized persistence modules)

Let $H : \mathbf{Top} \rightarrow \mathbf{A}$. Then

$$d(HF, HG) \leq d_\infty(f, g).$$

Examples

<i>persistence</i>	M	\mathbf{P}
<i>ordinary</i>	\mathbb{R}	$\{(-\infty, a] \mid a \in \mathbb{R}\}$
<i>multidimensional</i>	\mathbb{R}^n	$\{\mathbb{R}_{\leq a}^n \mid a \in \mathbb{R}^n\}$
<i>levelset</i>	\mathbb{R}	{intervals in \mathbb{R} }
<i>angle-valued</i>	S^1	{arcs in S^1 }
<i>cosheaf</i>	M	{open sets in M }

For each of these examples,

- \mathbf{P} is a poset under inclusion;
- for $f : X \rightarrow M$, $F : \mathbf{P} \rightarrow \mathbf{Top}$ is given by inverse images of f ;
- for $f, g : X \rightarrow M$ and $H : \mathbf{Top} \rightarrow \mathbf{A}$, $d(HF, HG) \leq d_\infty(f, g)$.

Comparing diagrams

A **natural transformation** is a map of diagrams.

For $V, W : \mathbf{n} \rightarrow \mathbf{Vect}_{\mathbb{F}}$ it is a commutative diagram,

$$\begin{array}{ccccccc}
 V_0 & \longrightarrow & V_1 & \longrightarrow & \cdots & \longrightarrow & V_n \\
 \varphi_0 \downarrow & & \varphi_1 \downarrow & & & & \downarrow \varphi_n \\
 W_0 & \longrightarrow & W_1 & \longrightarrow & \cdots & \longrightarrow & W_n
 \end{array}$$

For $F, G : \mathbf{P} \rightarrow \mathbf{D}$, for all $x \leq y$ there is a commuting diagram,

$$\begin{array}{ccc}
 F(x) & \xrightarrow{F(x \leq y)} & F(y) \\
 \varphi_x \downarrow & & \downarrow \varphi_y \\
 G(x) & \xrightarrow{G(x \leq y)} & G(y)
 \end{array}$$

Comparing diagrams

We denote a natural transformation by $\varphi : F \Rightarrow G$.

Two diagrams F, G are **isomorphic** if we have $\varphi : F \Rightarrow G$ and $\psi : G \Rightarrow F$ such that $\psi \circ \varphi = \text{Id}$ and $\varphi \circ \psi = \text{Id}$.

What if F and G are not isomorphic?

Comparing diagrams

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What if F and G are not isomorphic?

We would like to be able to quantify how far F and G are from being isomorphic.

- We will define **translations** on \mathbf{P} , and use these
- to define **interleavings** between diagrams.
- Then a metric on \mathbf{P} will give us the **interleaving distance**.

Translations

Definition

A **translation** is given by $\Gamma : \mathbf{P} \rightarrow \mathbf{P}$ such that $x \leq \Gamma(x)$ for all x .

- The identity is a translation.
- The composition of translations is a translation.

Interleaving

Definition

F and G are (Γ, K) -interleaved if there exist φ, ψ ,

$$\begin{array}{ccccc}
 \mathbf{P} & \xrightarrow{\Gamma} & \mathbf{P} & \xrightarrow{K} & \mathbf{P} \\
 \downarrow F & \xRightarrow{\varphi} & \downarrow G & \xRightarrow{\psi} & \downarrow F \\
 \mathbf{C} & \xlongequal{\quad} & \mathbf{C} & \xlongequal{\quad} & \mathbf{C}
 \end{array}$$

such that

$$\psi\varphi = FK\Gamma \text{ and } \varphi\psi = G\Gamma K.$$

Interleaving of (\mathbb{R}, \leq) -indexed diagrams

$$\Gamma = K : a \mapsto a + \varepsilon$$

$$\forall a \leq b,$$

$$\begin{array}{ccc}
 F(a) & \longrightarrow & F(b) \\
 \searrow & & \searrow \\
 & G(a + \varepsilon) \longrightarrow & G(b + \varepsilon)
 \end{array}
 \qquad
 \begin{array}{ccc}
 & F(a + \varepsilon) \longrightarrow & F(b + \varepsilon) \\
 & \nearrow & \nearrow \\
 G(a) & \longrightarrow & G(b)
 \end{array}$$

$$\forall a,$$

$$\begin{array}{ccc}
 F(a) & \longrightarrow & F(a + 2\varepsilon) \\
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 \nearrow & & \searrow \\
 G(a) & \longrightarrow & G(a + 2\varepsilon)
 \end{array}$$

Interleaving distance

Now assume that \mathbf{P} has a metric d .

An ε -translation is a translation $\Gamma : \mathbf{P} \rightarrow \mathbf{P}$ such that $d(x, \Gamma(x)) \leq \varepsilon$ for all x .

Definition

$$d(F, G) = \inf(\varepsilon \mid F, G \text{ interleaved by } \varepsilon\text{-translations})$$

Theorem (Interleaving distance)

This interleaving distance is metric.

Examples

Given (M, d_M) , let \mathbf{P} be a subset of $\mathcal{P}(M)$ with partial order given by inclusion, and Hausdorff distance.

<i>persistence</i>	M	\mathbf{P}
<i>ordinary</i>	\mathbb{R}	$\{(-\infty, a] \mid a \in \mathbb{R}\}$
<i>multidimensional</i>	\mathbb{R}^n	$\{\mathbb{R}_{\leq a}^n \mid a \in \mathbb{R}^n\}$
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- Particular ε -translations, Γ_ε , are given by thickening by ε .
- Note that each poset \mathbf{P} is closed under these Γ_ε .

Using functoriality

Theorem (Stability of interleaving distance)

Let $F, G : \mathbf{P} \rightarrow \mathbf{C}$ and $H : \mathbf{C} \rightarrow \mathbf{A}$. Then,

$$d(H \circ F, H \circ G) \leq d(F, G).$$

Proof.

$$\begin{array}{ccccc}
 \mathbf{P} & \xrightarrow{\Gamma} & \mathbf{P} & \xrightarrow{K} & \mathbf{P} \\
 \downarrow F & \Downarrow \varphi & \downarrow G & \Downarrow \psi & \downarrow F \\
 \mathbf{C} & \xlongequal{\quad} & \mathbf{C} & \xlongequal{\quad} & \mathbf{C} \\
 \downarrow H & = & \downarrow H & = & \downarrow H \\
 \mathbf{A} & \xlongequal{\quad} & \mathbf{A} & \xlongequal{\quad} & \mathbf{A}
 \end{array}$$



Inverse-image stability

Theorem (Inverse-image stability)

Let $F, G : \mathbf{P} \rightarrow \mathbf{Top}$ correspond to $f, g : X \rightarrow (M, d_M)$.
Assume \mathbf{P} closed under Γ_ε for all ε .

$$d(F, G) \leq d_\infty(f, g) := \sup_{x \in X} d_M(f(x), g(x)).$$

Proof.

Let $\varepsilon = d_\infty(f, g)$. F, G are ε -interleaved:

$$F(S) = f^{-1}(S) \subseteq g^{-1}(\Gamma_\varepsilon(S)) = G\Gamma_\varepsilon(S).$$

Thus, $d(F, G) \leq \varepsilon = d_\infty(f, g)$. □

Algebraic structure

Until this point we have only imposed structure on the indexing category: a partial order and a metric.

To compute we need some algebraic structure in our target category. For example, $\mathbf{Vect}_{\mathbb{F}_2}$, $\mathbf{Vect}_{\mathbb{F}}$, or \mathbf{Ab} .

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We will assume the target category, \mathbf{A} , is **abelian**.

This also includes **R-mod** and sheaves of abelian groups on X .

Definition

A category, \mathbf{A} , is **abelian** if

- each hom-set is an abelian group;
- all finite direct sums exist; and
- all morphisms have kernels and cokernels.

Kernel, image and cokernel persistence

Given ,

- $X \subseteq Y \in \mathbf{Top}$ and $g : Y \rightarrow (M, d_M)$;
- \mathbf{P} is a poset of subsets of M closed under Γ_ε , and
- $H : \mathbf{Top} \rightarrow \mathbf{A}$.

Let $f : X \hookrightarrow Y \xrightarrow{g} (M, d_M)$.

Let $F, G : \mathbf{P} \rightarrow \mathbf{Top}$ be given by inverse images of f, g .

Since $f^{-1}(U) \subset g^{-1}(U)$, $F \hookrightarrow G$, and $HF \xrightarrow{\alpha} HG \in \mathbf{A}^{\mathbf{P}}$.

Since \mathbf{A} is abelian, so is $\mathbf{A}^{\mathbf{P}}$.

So the kernel, image and cokernel of α exist.

Stability for kernel, image and cokernel persistence

Given $X \subseteq Y \in \mathbf{Top}$ and $g, g' : Y \rightarrow (M, d_M)$.

Construct $\alpha : HF \rightarrow HG$ and $\alpha' : HF' \rightarrow HG'$ as above.

Theorem (Stability of ker/im/coker persistence)

$$d(\ker(\alpha), \ker(\alpha')) \leq d_\infty(g, g')$$

$$d(\operatorname{im}(\alpha), \operatorname{im}(\alpha')) \leq d_\infty(g, g')$$

$$d(\operatorname{coker}(\alpha), \operatorname{coker}(\alpha')) \leq d_\infty(g, g')$$

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Monoid of translations

Recall def of translation: $\Gamma : \mathbf{P} \rightarrow \mathbf{P}$ such that $x \leq \Gamma(x)$ for all x .

That is, Γ is an endofunctor of \mathbf{P} together with $\text{Id} \Rightarrow \Gamma$.

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$\mathbf{End}_*(\mathbf{P})$ contain Id and are closed under composition.
So they are a monoid.

$$\mathbf{End}_*(\mathbf{P}) \begin{array}{c} \xrightarrow{\rho} \\ \xleftarrow{\iota} \end{array} (([0, \infty), \geq), +, 0)$$

$$\rho : \Gamma \mapsto \sup(d(x, \Gamma(x)))$$

$$\iota : \varepsilon \mapsto \Gamma_\varepsilon$$

Either of ρ or ι allows us to define interleaving distance.

Lawvere metric spaces

Instead starting with a poset \mathbf{P} and a metric, it suffices to have set with a function $d : \mathbf{P} \times \mathbf{P} \rightarrow [0, \infty]$ such that

- $d(a, a) = 0$ for all $a \in \mathbf{P}$, and
- d satisfies the triangle inequality.

That is, \mathbf{P} is an **Lawvere metric space**.

We define $a \leq b \in \mathbf{P}$ iff $d(a, b) < \infty$.

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A Lawvere metric space \mathbf{P} is a small category enriched over the monoidal poset $(([0, \infty], \geq), +, 0)$.

A preordered set is a small category enriched over the Boolean algebra $\{\text{True}, \text{False}\}$.

Our functor $\mathbf{Lawv} \rightarrow \mathbf{Proset}$ induced by the monoidal map $[0, \infty] \rightarrow \{\text{True}, \text{False}\}$ that maps ∞ to False and all else to True.

Main theorem

Our main results (Interleaving distance and stability of interleaving distance) can be summarized as follows.

Theorem

Given a poset \mathbf{P} with a metric, the interleaving distance gives a functor

$$\mathbf{Cat}(\mathbf{P}, -) : \mathbf{Cat} \rightarrow \mathbf{Metric}.$$